

POLYNOMIALS

Tutte polynomial

- e a bridge $\Leftrightarrow k(G-e) > k(G)$
- e a loop \Leftrightarrow incident to exactly one vertex

• Tutte polynomial is a polynomial valued graph invariant

• Tutte polynomial (1954)

$$T(G; x, y) \in \mathbb{Z}[x, y]$$

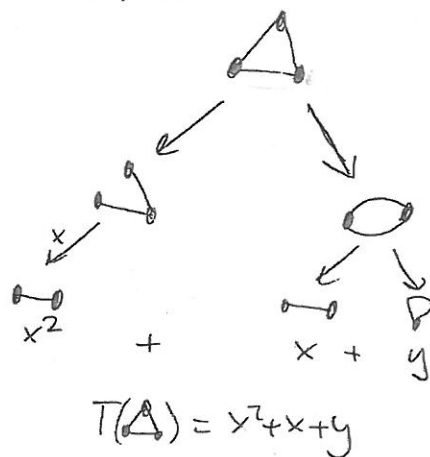
$$* T(G) = \begin{cases} x T(G/e) & e \text{ bridge} \\ y T(G/e) & e \text{ loop} \\ T(G \setminus e) + T(G/e) & e \text{ ordinary} \end{cases}$$

* $T(\cdot) = 1$

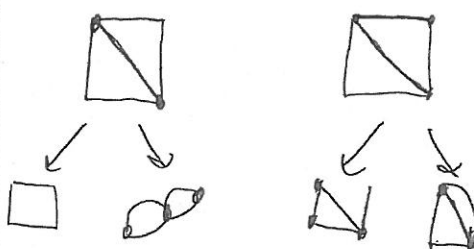
* $T(G \cup H) = T(G) \cdot T(H) = T(G \cup H)$

* $T(\text{two loops}) = T(\text{loop}) \cdot T(\text{loop}) = T(\text{two loops})$

Example



• There is a choice of order of resolution of edges:



• Thm: (i) $T(G)$ is well-defined

$$(ii) T(G) = \sum_{A \subseteq E(G)} (x-1)^{r(G)-r(A)} (y-1)^{n(A)}$$

• Tutte poly encodes lots of combinatorial info about a graph

* $(-1)^{|E|} \lambda^{k(G)} T(G; 1-\lambda, 0) = \text{chromatic poly} = \# \lambda\text{-cols of graph}$

* $(-1)^{n(G)} T(G; 0, 1-\lambda) = \# \text{ nowhere zero } \lambda\text{-flows}$

* $T(G; x, \frac{1}{x}) \rightsquigarrow$ Jones / HOMFLY poly

* $q^{k(G)} v^{r(G)} T(G; \frac{q}{v} + 1, v) \rightsquigarrow q$ state poly. $v = e^{2\pi i/n}$

* $T(2, 0) = \# \text{ acyclic orientations}$

* $T(1, 0) = \# \text{ acyclic orientations w/ one source}$

* $T(0, 2) = \# \text{ totally cyclic orientations (every edge in one cycle)}$

etc:

(lots of work: Handbook - 3 chapters)

Universality

- The Tutte poly is the most general deletion-contraction invariant for graphs, in the following sense.

Thm (Oxley, Welsh 79)

$\exists!$ map $U: \frac{\text{Graphs}}{is} \rightarrow \mathbb{Z}[x, y, \sigma, \tau]$

st.

$$U(v, \emptyset) = 1$$

$$U(G) = \begin{cases} x U(G/e) & e \text{ bridge} \\ y U(G/e) & e \text{ loop} \\ \sigma U(G/e) + \tau U(G/e) & \text{otherwise} \end{cases}$$

Moreover

$$U(G) = \sigma^{n(G)} \tau^{r(G)} T(G; \frac{x}{\tau}, \frac{y}{\sigma})$$

Ribbon graphs

- What is the "Tutte polynomial" of a ribbon graph?
- What do we mean by Tutte poly?
 - * universal del-cont invariant
 - * have deletion and contraction.
- Need to know cases. For this examine universality for $T(G)$

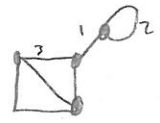
- recognising bridges, loops and ordinary edges via deletion contraction.

$$e^c := E \setminus e$$

$$G/e^c = \begin{cases} \text{---} & \because \text{bridge} = b \\ \bigcirc & \because \text{loop} = l \end{cases}$$

$$G/e^c = \begin{cases} \text{---} & \because \\ \bigcirc & \because \end{cases}$$

- Say for $i, j \in \{l, b\}$
 e is (i, j) -edge $\Leftrightarrow \begin{cases} G/e^c \text{ is } i \\ G \times e^c \text{ is } j \end{cases}$

example: 

- 1 is (b, b)
- 2 is (l, l)
- 3 is (b, b)

- (b, b) \Leftrightarrow bridge
- (l, l) \Leftrightarrow loop
- (l, b) otherwise
- (b, l) impossible

- variables:
 - a_i when G/e^c type i
 - b_i when $G \times e^c$ type j

- Universality becomes.

$\exists!$ U st

$$U(v, \emptyset) = 1$$

$$U(G) = a_i U(G/e) + b_j U(G/e) \quad \text{type } (i, j)$$

moreover

$$U(G) = \sigma^{n(G)} \tau^{r(G)} T(G; \frac{a_b}{b_b} + 1, \frac{b_l}{a_l} + 1)$$

Applying this to ribbon graphs

* $G/e^c, G/e^c \in \left\{ \begin{matrix} \text{---} \\ b \end{matrix}, \begin{matrix} \text{---} \\ 0 \end{matrix}, \begin{matrix} \text{---} \\ n \end{matrix} \right\}$

* $i, j \in \{b, 0, n\}$

$e \text{ is } (i, j) \Leftrightarrow \begin{cases} G/e^c \text{ is } i \\ G/e^c \text{ is } j \end{cases}$

example

$G/1^c = \text{---}, G/1^c = \text{---} \text{ (no)}$
 $G/2^c = \text{---}, G/2^c = \text{---} \text{ (no)}$

variables: a_i when G/e^c type i
 b_i when G/e^c type j

Theorem:

$\exists!$ invariant α satisfying

$\alpha(\text{edgless}) = 1$

$\alpha(G) = a_i \alpha(G/e) + b_j \alpha(G/e)$
if $e \text{ is } (i, j)$

if and only if

$a_n = \sqrt{a_b a_0}, b_n = \sqrt{b_b b_0}$

in which case

$\alpha = a_0^{|\mathcal{E}| - \ell(G)} b_b^{\ell(G)}$
 $\sum_{A \subseteq \mathcal{E}} \left(\frac{a_b}{b_0} \right)^{\ell(G) - \ell(A)} \left(\frac{b_b}{a_0} \right)^{|\mathcal{A}| - \ell(A)}$

where

$\ell(A) := \frac{1}{2} (|\mathcal{A}| + v(A) - f(A))$
 $= r(A) + \frac{1}{2} \delta(A)$

so Define " Tutte poly " of a ribbon graph by

$\tilde{R}(G; x, y) := \sum_{A \subseteq \mathcal{E}} (x-1)^{\ell(G) - \ell(A)} (y-1)^{|\mathcal{A}| - \ell(A)}$

Found in earlier lit as follows

Pollak's - Norden $(0, 0, 2)$

$R(G; x, y, z) = \sum (x-1)^{\ell(G) - \ell(A)} y^{|\mathcal{A}| - \ell(A)} z^{\delta(A)}$

$\tilde{R}(G; x+1, y+1) = x^{\frac{1}{2}\delta(G)} R(G; x+1, y, \sqrt{xy})$

[Remarks on this poly in lit.]

proof:

• computing $\alpha(\text{---})$ in two ways gives the condition on the variables.

• Show state sum satisfies def-conds.

by splitting sum $\sum_{A \subseteq \mathcal{E}} \sum_{A \subseteq \mathcal{E}}$

rewriting in terms of $\alpha(G/e)$ & $\alpha(G)$

□

